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## ON THE HORTON-STRAHLER NUMBER FOR COMBINATORIAL TRIES

MARKUS E. NEBEL<sup>1</sup>

**Abstract.** In this paper we investigate the average Horton-Strahler number of all possible tree-structures of binary tries. For that purpose we consider a generalization of extended binary trees where leaves are distinguished in order to represent the location of keys within a corresponding trie. Assuming a uniform distribution for those trees we prove that the expected Horton-Strahler number of a tree with  $\alpha$  internal nodes and  $\beta$  leaves that correspond to a key is asymptotically given by

$$\frac{4^{2\beta-\alpha} \log(\alpha)(2\beta-1)(\alpha+1)(\alpha+2) \binom{2\alpha+1}{\alpha-1}}{8\sqrt{\pi}\alpha^{3/2} \log(2)(\beta-1)\beta \binom{2\beta}{\beta}^2}$$

provided that  $\alpha$  and  $\beta$  grow in some fixed proportion  $\rho$  when  $\alpha \rightarrow \infty$ . A similar result is shown for trees with  $\alpha$  internal nodes but with an arbitrary number of keys.

**AMS Subject Classification.** 05A15, 05C05, 68W40.

### 1. INTRODUCTION

Let  $T$  be a binary tree, *i.e.* a tree where each node has at most two descendants. Then the *Horton-Strahler number* of  $T$  denoted  $hs(T)$  is recursively defined by

$$hs(T) := \begin{cases} 0 & : T \text{ is either a leaf or empty} \\ hs(T.l) + 1 & : \text{if } hs(T.l) = hs(T.r) \\ \max(hs(T.l), hs(T.r)) & : \text{otherwise.} \end{cases}$$

Here,  $T.l$  (resp.  $T.r$ ) denotes the left (resp. right) subtree of  $T$ . The Horton-Strahler number was originally introduced to classify river systems (see [12] and [22]) but it has also been adopted in computer science, molecular biology, medicine and other disciplines. Ershov [6], for example, has shown that the minimal

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number of registers needed to evaluate an arithmetic expression  $\mathcal{E}$  with binary operators, which is represented as a binary tree  $T(\mathcal{E})$  (the syntax-tree), is given by  $1 + hs(T(\mathcal{E}))$ . If all syntax-trees with  $n$  internal nodes ( $n$  binary operators) are assumed to be equally likely, then it is known that the average number of registers that are needed to evaluate an expression optimally is given by

$$\frac{1}{2} \log_2(2\pi^2 n) - \frac{\gamma + 2}{2 \ln(2)} + F(n) + \mathcal{O}(n^{-1/2+\delta}) \quad (1)$$

for all  $\delta > 0$ ,  $n \rightarrow \infty$ , where  $\gamma = 0.577215\dots$  is Euler's constant and  $F$  is a periodic, oscillating function of small amplitude (see *e.g.* [7] and [13]). Syntax-trees corresponding to expressions built with unary and binary operators were considered in [9]. Furthermore, the minimum stack-size required for a traversal of a binary tree  $T$  is also given by  $1 + hs(T)$  (*e.g.* see [7] and [11]). Meir *et al.* [18] investigated the Horton-Strahler number of channel networks with a fixed number of inputs. The combinatorics of the Horton-Strahler analysis has been used in computer graphics for the creation of faithful synthetic images of trees (see [23]). The impact of the Horton-Strahler number on molecular biology comes from theoretical considerations about secondary structures of single-stranded nucleic acids (see [24] and the references given there).

All those applications and studies have in common that they deal with ordinary extended binary trees, *i.e.* trees where each node is either a leaf or has two descendants. All the cited papers which present an average case analysis consider the uniform model, *i.e.* they assume that all trees of a given size are equally likely.

The Horton-Strahler number of tries has been investigated in a recent work by Devroye and Kruszewski [4]. A trie is a binary tree which is used to store the set of keys  $K = \{k_1, \dots, k_n\}$  in the following manner: each key  $k_i$ , considered as a string of 0's and 1's due to its binary representation, defines a path in a binary tree  $T$  (0 indicates a left turn, 1 a right turn); the trie defined by  $k_1, \dots, k_n$  is the smallest binary tree for which the paths truncated at the leaves of  $T$  are all pairwise different. Thus each leaf of  $T$  stores exactly one of the keys  $k_i$ ,  $1 \leq i \leq n$ . Note that  $T$  does not need to be an extended binary tree.  $T$  might have internal nodes with only one successor. However, the Horton-Strahler number  $hs(T)$  remains unchanged when we turn  $T$  into an extended binary tree. Devroye and Kruszewski considered random tries constructed from  $n$  i.i.d. sequences of Bernoulli random variables with parameter  $p$ ,  $0 < p < 1$ ; they have shown that the Horton-Strahler number  $H_n$  of those tries fulfils

$$\frac{H_n}{\log n} \rightarrow \frac{1}{\log \frac{1}{\min(p, 1-p)}}$$

in probability as  $n \rightarrow \infty$ . The presented (Bernoulli-) model of a random trie is very realistic. For example, if we choose  $p$  being  $\frac{1}{2}$ , this model describes exactly the behavior of tries built from random integer data assuming all integers to be equally likely.



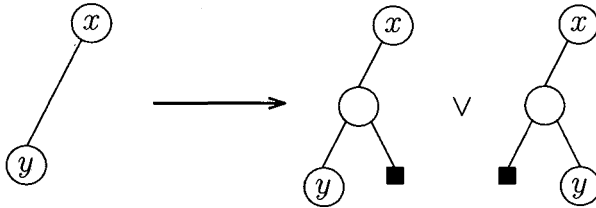


FIGURE 2. Affecting additive parameters by adding a new node in between two existing ones. All trees shown in the figure have the same Horton-Strahler number but different path lengths.

length of a river should not affect the classification of a river network). However, we change length-sensitive parameters like *e.g.* the external path length.

Thus, even if all our results are presented using the term  $\mathcal{C}$ -trie instead of generalized extended binary tree they should be considered as applicable to many other areas.

## 2. THE AVERAGE HORTON-STRAHLER NUMBER

The aim of this section is to derive the average Horton-Strahler number for uniform random  $\mathcal{C}$ -tries as defined in Section 1, *i.e.* the average value of the function  $hs$  applied to a set of  $\mathcal{C}$ -tries of the same size  $(\alpha, \beta)$ . We will use generating functions in order to prove our results. The way these generating functions are derived is similar to that in [7], the methodology used to determine asymptotics for the coefficients in question is standard and can be found in [8]. By  $[x_1^{n_1} \dots x_k^{n_k}]f(x_1, \dots, x_k)$  we denote the coefficient at  $x_1^{n_1} \dots x_k^{n_k}$  in an expansion of  $f(x_1, \dots, x_k)$  at  $(x_1, \dots, x_k) = (0, \dots, 0)$ .

We start our investigations by determining the generating function  $H_p(x, y)$  which counts those  $\mathcal{C}$ -tries that have a Horton-Strahler number of exactly  $p$ .

**Lemma 1.** *Let  $x$  mark an internal node and let  $y$  mark a white leaf. The generating function  $H_p(x, y)$  of  $\mathcal{C}$ -tries  $T$  with  $hs(T) = p$  possesses the following closed form representation:*

$$H_p(x, y) = \frac{\sin(\phi)}{\sin(2^{p-1}\phi)} \cdot \frac{xy^2}{1 - 2x - 2xy}, \tag{2}$$

$$\text{where } \phi = \arccos\left(\frac{1 - 4x(y + 1) + 2x^2(2 + y(4 + y))}{2x^2y^2}\right). \tag{3}$$

*Proof.* In order to derive a representation for the generating function in question we have to distinguish the cases shown in Figure 3. For  $p \geq 2$  these cases translate

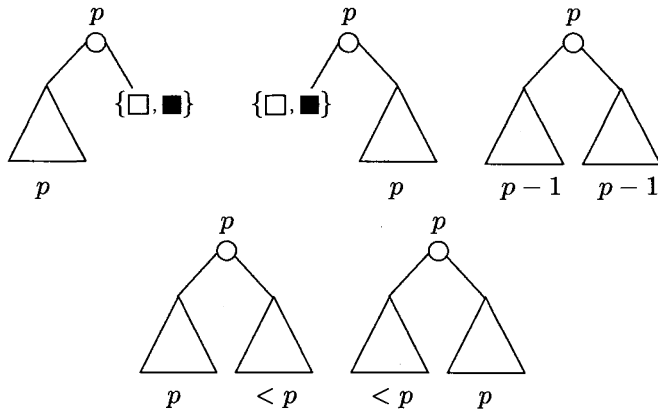


FIGURE 3. All possible decompositions of a  $\mathcal{C}$ -trie  $T$  with  $hs(T) = p$ . The number below a triangle determines the value of  $hs$  of the subtree represented.

into the following recurrence for  $H_p(x, y)$ :

$$H_p(x, y) = (2x + 2xy)H_p(x, y) + xH_{p-1}^2(x, y) + 2xH_p(x, y) \sum_{1 \leq j < p} H_j(x, y). \quad (4)$$

Here, the boundary condition for  $p = 1$  is still to be determined yet. It is obvious that a  $\mathcal{C}$ -trie  $T$  with  $hs(T) = 1$  has to have a linear structure, *i.e.* either the left or the right subtree of each internal node has to be a leaf. Thus  $H_1(x, y)$  must fulfil  $H_1(x, y) = xy^2 + (2x + 2xy)H_1(x, y)$  and therefore

$$H_1(x, y) = \frac{xy^2}{1 - 2x - 2xy}.$$

In order to solve this recurrence we divide both sides of (4) by  $xH_p(x, y)$ , thus

$$\frac{1}{x} = 2 + 2y + \frac{H_{p-1}^2(x, y)}{H_p(x, y)} + 2 \sum_{1 \leq j < p} H_j(x, y).$$

Subtracting this from the analogous identity obtained for  $p + 1$  eliminates the summation. We find

$$0 = \frac{H_p^2(x, y)}{H_{p+1}(x, y)} + 2H_p(x, y) - \frac{H_{p-1}^2(x, y)}{H_p(x, y)}. \quad (5)$$

Let  $V_p(x, y) := \frac{H_{p-1}(x, y)}{H_p(x, y)}$ . Dividing (5) by  $H_p(x, y)$  our recurrence can be expressed by means of  $V_p$ :

$$V_{p+1}(x, y) = V_p^2(x, y) - 2, \quad p \geq 2,$$

with the initial condition  $V_2(x, y) = \frac{H_1(x, y)}{H_2(x, y)}$ . We can determine  $H_2(x, y)$  by using  $H_1(x, y)$  and the recurrence (4) which yields

$$V_2(x, y) = \frac{1 - 4x(y + 1) + 2x^2(2 + y(4 + y))}{x^2y^2}.$$

This new recurrence can be solved by a trigonometric change of variables. We set  $V_2(x, y) = 2 \cos(\phi)$  and generally  $V_p(x, y) = 2 \cos(\phi_p)$ . Since  $\cos^2(x) = \frac{1}{2}(1 + \cos(2x))$  holds we see that the recurrence translates into

$$2 \cos(\phi_{p+1}) = 2 \cos(2\phi_p).$$

Therefore, for  $p \geq 2$ ,  $\phi_{p+1} = 2\phi_p = 2^{p-1}\phi$  must hold which gives the explicit form

$$\begin{aligned} V_{p+1}(x, y) &= 2 \cos(2^{p-1}\phi), \quad p \geq 2, \\ V_2(x, y) &= 2 \cos(\phi) = \frac{1 - 4x(y + 1) + 2x^2(2 + y(4 + y))}{x^2y^2}. \end{aligned}$$

We can go back to  $H_p(x, y)$  by regarding

$$V_p(x, y)V_{p-1}(x, y) \cdots V_2(x) = \frac{H_{p-1}(x, y)}{H_p(x, y)} \frac{H_{p-2}(x, y)}{H_{p-1}(x, y)} \cdots \frac{H_1(x, y)}{H_2(x, y)} = \frac{H_1(x, y)}{H_p(x, y)}.$$

By means of the identity  $\sin(2x) = 2 \sin(x) \cos(x)$  the product on the left-hand side collapses to  $\sin(2^{p-1}\phi)$  when multiplied by  $\sin(\phi)$ . This completes the proof. □

Next we consider those  $\mathcal{C}$ -tries that have a Horton-Strahler number of at least  $p$ .

**Lemma 2.** Let  $S_p(x, y) := \sum_{j \geq p} H_j(x, y)$ ,  $\kappa := \frac{x^2y^2}{(1-2x-2xy)^2}$ ,  $\varepsilon := \sqrt{1-4\kappa}$  and  $u := \frac{1-\varepsilon}{1+\varepsilon}$ . We have

$$S_p(x, y) = y \cdot \left[ \frac{\sqrt{u}(1-u)}{u} \cdot \frac{u^{2^{p-1}}}{1-u^{2^{p-1}}} \right].$$

*Proof.* In order to prove the lemma we use the identity  $\sin(x) = \frac{1}{2i}(e^{ix} - e^{-ix})$ ,  $i^2 = -1$ , which we insert in (2). Together with  $t := e^{-i\phi}$  and  $r := 2^{p-1}$  we find

$$H_p(x, y) = 2i \sin(\phi) \frac{t^r}{1-t^{2r}} \frac{xy^2}{1-2x-2xy}.$$

Now consider those parts of the representation that depend on  $p$ . Summing them up for  $j \geq p$  yields

$$\sum_{\substack{j \geq p \\ r=2^{j-1}}} \frac{t^r}{1-t^{2r}} = \sum_{\substack{j \geq p \\ \kappa \geq 0}} t^{2^{j-1}(1+2\kappa)} = \sum_{\substack{m, \kappa \geq 0 \\ v=2^{p-1}}} v^{2^m(2\kappa+1)}.$$

Since the mapping  $(m, k) \rightarrow 2^m(2k + 1)$  is a bijection on  $\mathbb{N}^2 \rightarrow \mathbb{N}$  the last sum equals  $\frac{v}{1-v}$ . Therefore  $S_p(x, y) = 2i \sin(\phi) \frac{xy^2}{1-2x-2xy} \frac{t^r}{1-t^r}$  holds. Returning to trigonometric functions, *i.e.* setting  $t = e^{-i\phi} = \cos(\phi) - i \sin(\phi)$ , gives us

$$S_p(x, y) = \left[ -i \sin(\phi) + \frac{\sin(\phi) \cos(2^{p-2}\phi)}{\sin(2^{p-2}\phi)} \right] \frac{xy^2}{1-2x-2xy}.$$

For  $p > 1$  it is now possible to express the trigonometric functions by means of Chebyshev polynomials. For  $p = 1$  we run into trouble since in that case we would refer to the  $\frac{1}{2}$ -th polynomial which does not exist. Thus, the next step is to express  $S_p(x, y)$  by means of  $\frac{1}{2}\phi$  instead of  $\phi$ . By applying the identity  $2 \cos(x) \sin(x) = \sin(2x)$  we find that

$$S_p(x, y) = \frac{-i \sin(\phi)xy^2}{1-2x-2xy} + \frac{\cos(\frac{1}{2}\phi) \sin(\frac{1}{2}\phi) \cos(2^{p-1}\frac{1}{2}\phi)}{\sin(2^{p-1}\frac{1}{2}\phi)} \frac{2xy^2}{1-2x-2xy}$$

holds. From equation (3) we derive closed form expressions for  $i \sin(\phi)$  and  $\cos(\frac{1}{2}\phi)$  which we insert into the last representation of  $S_p(x, y)$ . Then, applying the following identities for the Chebyshev polynomial of the first-kind ( $T_n(x)$ ) (see *e.g.* [1] 22.3.15) and the second-kind ( $U_n(x)$ ) (see *e.g.* [1] 22.3.16)

$$T_n(\cos(\phi)) = \cos(n\phi),$$

$$U_n(\cos(\phi)) = \frac{\sin((n+1)\phi)}{\sin(\phi)},$$

yield

$$S_p(x, y) = -\frac{\sqrt{(1-2x)(1-2x-4xy)}}{2x} - \frac{yT_{2^{p-1}}(\hat{\kappa})}{U_{2^{p-1}-1}(\hat{\kappa})}.$$

Here  $\hat{\kappa} := \cos(\frac{1}{2}\phi) = -\frac{1-2x-2xy}{2xy}$  holds. Now let  $T(x, y)$  be the ordinary generating function of all  $\mathcal{C}$ -tries. In [21] the following representation can be found:

$$T(x, y) = \frac{1-2x-\sqrt{(1-2x)(1-2x-4xy)}}{2x}.$$

This, together with two further identities for Chebyshev polynomials

$$T_n(x) = U_n(x) - xU_{n-1}(x), \text{ (see e.g. [1], 22.5.6),}$$

$$U_{n+1}(x) = 2xU_n(x) - U_{n-1}(x), \text{ (see e.g. [15] (B77)),}$$

gives us

$$S_p(x, y) = T(x, y) - y + \frac{yU_{2^{p-1}-2}(\hat{\kappa})}{U_{2^{p-1}-1}(\hat{\kappa})},$$



for  $U_{-1}(x) := 0$ . A closed form representation for  $U_n(x)$  is given in [15] (B74). By fundamental algebraic manipulations this representation can be transformed into

$$U_n(x) = -\frac{x^n [(1 - \sqrt{1 - x^{-2}})^{n+1} - (1 + \sqrt{1 - x^{-2}})^{n+1}]}{2\sqrt{1 - x^{-2}}}.$$

Now, since  $1 - \hat{\kappa}^{-2} = 1 - 4\kappa$  holds, we get

$$S_p(x, y) = T(x, y) - y - 2y\sqrt{\kappa} \frac{(1 + \sqrt{1 - 4\kappa})^{2^{p-1}-1} - (1 - \sqrt{1 - 4\kappa})^{2^{p-1}-1}}{(1 + \sqrt{1 - 4\kappa})^{2^{p-1}} - (1 - \sqrt{1 - 4\kappa})^{2^{p-1}}}.$$

We complete the proof by using the substitutions of the lemma and applying some obvious simplifications. □

**Remark.** Besides the Horton-Strahler number there is another monotonic marking of binary trees which is related to the evaluation of arithmetic expressions and the traversal. This is the so called *stack-number* of the tree which was investigated in numerous papers (e.g. [3, 14, 16, 19, 20] and [21]). It corresponds to the stack-size needed to traverse a tree in preorder using the traditional algorithm (see e.g. [17], p. 319ff) and the number of cells on a stack needed to evaluate an arithmetic expression by means of a simple traversal algorithm (see [15] for details). For non-colored extended binary trees we have the following correspondence: the number of trees with  $\alpha$  internal nodes, with a stack-number of at most  $2^k - 1$ , is equal to the number of trees with  $\alpha$  internal nodes and a Horton-Strahler number of  $k$  (see e.g. [15], Th. 5.8). If we inspect the generating functions of the previous lemma and of [20] and [21] we see that such a relation does not exist for  $\mathcal{C}$ -tries neither of size  $\alpha$  nor of size  $(\alpha, \beta)$ .

In order to compute the average Horton-Strahler number we need a representation of the generating function  $M(x, y) := \sum_{p \geq 1} pH_p(x, y)$ . It is not hard to see that  $M(x, y) = \sum_{p \geq 1} S_p(x, y)$  holds. Therefore we have

$$\begin{aligned} M(x, y) &= \frac{y\sqrt{u}(1-u)}{u} \sum_{p \geq 1} \frac{u^{2^{p-1}}}{1-u^{2^{p-1}}} \\ &= \frac{y\sqrt{u}(1-u)}{u} \sum_{n \geq 1} (v_2(n) + 1)u^n. \end{aligned} \tag{6}$$

Here  $v_2(n)$  denotes the *dyadic valuation* of  $n$ , i.e. the number of positive divisors of  $n$  which are a power of two.

Now, everything is prepared to determine an asymptotic equivalent for the average Horton-Strahler number. We use the *Mellin summation method* as described in [10] to evaluate the sum. For that purpose we set  $u = \exp(-t)$  and

apply the well-known identity

$$\exp(-tj) = \frac{1}{2\pi i} \int_{c-i\infty}^{c+i\infty} \Gamma(s)j^{-s}t^{-s}ds, \quad i^2 = -1,$$

for some  $c$  in the fundamental strip of the Mellin transform of  $\exp(-tj)$  and for  $\Gamma(s)$  the complete gamma function. This is how it is possible to express the number-theoretic function  $v_2(n)$  by means of the Riemann Zeta function  $\zeta(z)$  (see [2] for details). We have (see e.g. [15], p. 155)

$$\sum_{n \geq 1} v_2(n)n^{-z} = \sum_{n \geq 1} v_2(2n)(2n)^{-z} = \sum_{\substack{j \geq 1 \\ n \geq 1}} (2^j n)^{-z} = \zeta(z)(2^z - 1)^{-1}$$

and therefore with  $u = \exp(-t)$

$$\sum_{n \geq 1} (v_2(n) + 1)u^n = \frac{1}{2\pi i} \int_{c-i\infty}^{c+i\infty} \zeta(s)2^s\Gamma(s)t^{-s}(2^s - 1)^{-1}ds.$$

Now, according to the Mellin summation formula, we have to sum the residues of  $\zeta(s)2^s\Gamma(s)t^{-s}(2^s - 1)^{-1}$  left to the fundamental strip, *i.e.* the residues with a real part less or equal to one. There are singularities at  $s = 1$  and  $s = -n$ ,  $n \in \mathbb{N} \cup \{0\}$ , but we only have to consider those which are larger than  $-1$  since the others will only imply terms that can be neglected. There are further singularities at  $s = \frac{2\pi ik}{\ln(2)} =: \chi_k$ ,  $k \in \mathbb{Z} \setminus \{0\}$ , which would imply some oscillation in the lower order terms. As the known methods for multivariate asymptotics only allow to determine the leading term, the singularities  $\chi_k$  will only be considered later in the univariate case. The sum of the residues for  $s = 1$  and  $s = 0$  is given by

$$\frac{2}{t} + \frac{2 \ln(t) + 2\gamma - 2 \ln(\pi) - 3 \ln(2)}{4 \ln(2)}.$$

Here,  $\gamma = 0.5772156649\dots$  denotes Eulers's constant.

In order to approximate the coefficient of  $M(x, y)$  at  $x^\alpha y^\beta$  we are interested in an expansion of our generating function at its dominant singularity. For that purpose we assume  $y$  being a positive constant not equal to 0 in order to determine the dominant singularity with respect to  $x$ , *i.e.* the value of  $x$  which has the smallest modulus and which is a singular point of  $M(x, y)$ . Note that this approach leads to the restriction that our asymptotic will only be valid when  $\alpha$  and  $\beta$  grow simultaneously in a fixed proportion. By definition of  $M(x, y)$  and properties of the Horton-Strahler number we have the following trivial bounds for  $[x^\alpha]M(x, y)$ :

$$[x^\alpha]T(x, y) \leq [x^\alpha]M(x, y) \leq \log_2(\alpha + 1)[x^\alpha]T(x, y).$$

Under the assumption that  $y$  is a constant and since  $x = \frac{1}{2+4y}$  is the dominant singularity of  $T(x, y)$  the  $\mathcal{O}$ -transfer method introduced in [8] leads to

$$[x^\alpha]T(x, y) \sim \frac{1 + 2y}{2} \cdot \frac{(2 + 4y)^\alpha}{\sqrt{\pi\alpha^3}}.$$

Therefore the Cauchy-Hadamard formula tells us that both the minorant and the majorant of  $M(x, y)$  have a radius of convergence of  $\frac{1}{2+4y}$  and we infer that  $\frac{1}{2+4y}$  is the radius of convergence of  $M(x, y)$  itself. Thus, by the theorem of Pringsheim, we can conclude that  $x = \frac{1}{2+4y}$  is a dominant singularity of our generating function. It resides to prove that it is the only dominant singularity. For that purpose we consider the representation (6) of  $M(x, y)$ . Besides the algebraic singularity at  $x = \frac{1}{2+4y}$  implied by our substitution, the factor of the sum is only singular for  $u = 0$  i.e. for  $x = 0$ . This is why it does not extend the set of dominant singularities. The sum  $\sum_{n \geq 1} (v_2(n)+1)u^n$  possesses the minorant  $\sum_{n \geq 1} u^n$  and the majorant  $\sum_{n \geq 1} (n+1)u^n$  both with a radius of convergence equal to 1. Therefore the set of solutions of  $|u| = 1$  might contribute further dominant singularities. However,  $|u| = 1$  has only one solution with an appropriate modulus, namely  $x = \frac{1}{2+4y}$ . Thus we can conclude that there is only one dominant singularity. Since  $t = -\log\left(\frac{1-\varepsilon}{1+\varepsilon}\right)$  and  $\varepsilon$  becomes 0 at our dominant singularity we expand  $-\log\left(\frac{1-\varepsilon}{1+\varepsilon}\right)$  about  $\varepsilon = 0$  to get  $t \sim 2\varepsilon$  and thus  $t \sim 2\frac{\sqrt{(1-2x)(1-2x-4xy)}}{1-2x-2xy}$ . We conclude that for an expansion at  $x = \frac{1}{2+4y}$ ,  $t$  complies with  $2\frac{\sqrt{2y(1+2y)}}{y}\sqrt{1-x(2+4y)}$ . On the assumption that  $y$  is constant and for  $u = \exp(-t)$ , the factor  $\frac{y\sqrt{u(1-u)}}{u}$  possesses the expansion  $yt + \mathcal{O}(t^2)$ . So, the most significant term of the expansion of  $M(x, y)$  around  $x = \frac{1}{2+4y}$  is given by

$$\frac{yt \ln(t)}{2 \ln(2)} \sim -\frac{\sqrt{2y(1+2y)}}{2 \log(2)} \sqrt{1-x(2+4y)} \log((1-x(2+4y))^{-1}). \tag{7}$$

This representation can be used to approximate the coefficients of  $M(x, y)$ . We find:

**Lemma 3.** *Let  $\rho := \frac{\alpha}{\beta}$  be fixed. The coefficient of  $M(x, y)$  at  $x^\alpha y^\beta$  is asymptotically given by*

$$\frac{2^{\alpha(1+\frac{1}{\rho})-2}}{\sqrt{\pi\alpha^3}} \left(\alpha + \frac{1}{2}\right) \log_2(\alpha),$$

$\alpha \rightarrow \infty$ .

*Proof.* We use the following well-known expansions

$$\begin{aligned} \log((1 - x(2 + 4y))^{-1}) &= \sum_{n \geq 1} \frac{x^n}{n} \sum_{k \geq 0} \binom{n}{k} 2^{n-k} 4^k y^k, \\ \sqrt{1 - x(2 + 4y)} &= \sum_{i \geq 0} \binom{\frac{1}{2}}{i} (-1)^i x^i \sum_{k \geq 0} \binom{i}{k} 2^{i-k} 4^k y^k, \\ \sqrt{2y + 4y^2} &= \sum_{j \geq 0} \binom{\frac{1}{2}}{j} 2^{1-j} y^{1-j}, \end{aligned}$$

and extract the coefficient at  $x^\alpha y^\beta$  in the resulting expansion of the right-hand side of (7). We find

$$[x^\alpha y^\beta]M(x, y) \sim -\frac{2^{\alpha+\beta-1}}{\log(2)} \binom{\alpha + \frac{1}{2}}{\beta - \frac{1}{2}} \underbrace{\sum_{n \geq 1} \frac{(-1)^{\alpha-n}}{n} \binom{\frac{1}{2}}{\alpha - n}}_{=: \sigma(\alpha)}.$$

By induction on  $\alpha$  it is possible to prove the following recursion for  $\sigma(\alpha)$ :

$$\begin{aligned} \sigma(0) &= 0, \\ \sigma(\alpha) &= \frac{2\alpha - 3}{2\alpha} \sigma(\alpha - 1) + \underbrace{\frac{(-1)^{\alpha+1} \sqrt{\pi}}{2\Gamma(\frac{5}{2} - \alpha)\Gamma(\alpha + 1)}}_{=: \varsigma(\alpha)}. \end{aligned}$$

This recursion can be solved by using ordinary generating functions. For  $A(z) := \sum_{\alpha \geq 0} \sigma(\alpha) z^\alpha$  we get  $A(z) = zA(z) - \frac{3}{2} \int_0^z A(t) dt + \sum_{\alpha \geq 1} \varsigma(\alpha) z^\alpha$ . Applying the identity  $\sum_{\alpha \geq 1} \varsigma(\alpha) z^\alpha = \frac{2}{3} \sqrt{1-z} z(z-1) + \frac{2}{3}$ , which for instance Zeilberger’s “fast algorithm” (see [25]) finds for you, yields a simple differential equation for  $A(z)$  which possesses the solution  $A(z) = -\sqrt{1-z} \log(1-z)$  and thus

$$\sigma(\alpha) = [z^\alpha] \sqrt{1-z} \log((1-z)^{-1})$$

holds. By applying the  $\mathcal{O}$ -transfer method we find the approximation  $\sigma(\alpha) \sim -\frac{\log(\alpha)}{2\sqrt{\pi\alpha^3}}$  which proves the lemma.  $\square$

In order to get the average value, the coefficient given in Lemma 3 has to be divided by the total number of  $\mathcal{C}$ -tries of size  $(\alpha, \beta)$ . We can proceed in the same way as done in the previous proof in order to approximate the coefficient  $[x^\alpha y^\beta]T(x, y)$ . Thus we factor  $T(x, y)$  into

$$\frac{1 - 2x}{2x} - \frac{1 - 2x}{2x} \sqrt{1 - 4xy(1 - 2x)^{-1}} \tag{8}$$

which corresponds to an expansion of  $T(x, y)$  around the singularity  $y = \frac{1-2x}{4x}$  (assuming now that  $x$  is constant). Since the leftmost term of (8) only possesses coefficients at  $y^0$  it can be neglected. Furthermore,  $\sqrt{1 - 4xy(1 - 2x)^{-1}}$  can be expanded using the binomial theorem, which yields

$$\xi(\alpha, \beta) := [x^\alpha y^\beta] \sqrt{1 - 4xy(1 - 2x)^{-1}} = \binom{\frac{1}{2}}{\beta} \binom{\alpha - 1}{\alpha - \beta} (-1)^\beta 4^\beta 2^{\alpha - \beta}.$$

Now, taking the factor  $-\frac{1-2x}{2x}$  into account proves that  $[x^\alpha y^\beta]T(x, y) \sim \xi(\alpha, \beta) - \frac{1}{2}\xi(\alpha + 1, \beta)$ . Thus we conclude that

$$[x^\alpha y^\beta]T(x, y) \sim \frac{2^{\alpha - \beta + 1}}{\beta} \binom{2\beta - 2}{\beta - 1} \binom{\alpha - 1}{\alpha - \beta + 1}.$$

Dividing the coefficient given in Lemma 3 by the previous quantity provides the following theorem:

**Theorem 1.** *On the assumption that all  $(\alpha, \beta)$ -tries are equally likely the average Horton-Strahler number of a  $\mathcal{C}$ -trie of size  $(\alpha, \beta)$  is asymptotically given by*

$$\frac{4^{\alpha(\frac{2}{\rho} - 1)} \log(\alpha) (2\frac{\alpha}{\rho} - 1) (\alpha + 1) (\alpha + 2) \binom{2\alpha + 1}{\alpha - 1}}{8\sqrt{\pi} \alpha^{3/2} \log(2) (\frac{\alpha}{\rho} - 1) \binom{\frac{2\alpha}{\rho}}{\frac{\alpha}{\rho}}^2},$$

$\rho := \frac{\alpha}{\beta}$  fixed,  $\alpha \rightarrow \infty$ .

**Remark.** The asymptotic given for the number of  $(\alpha, \beta)$ -tries is equal to the exact number of  $\mathcal{C}$ -tries of this size for  $\alpha > 0$ . This is due to the fact that we find a factorization of  $T(x, y)$  when expanding it around its dominant singularity. Thus, besides some terms at  $y^0$ , no terms were neglected when we have developed the leading term and have extracted the coefficients.

Looking at a plot of our average Horton-Strahler number of  $\mathcal{C}$ -tries (see the last section of this paper) it seems to be hardly dependent on  $\beta$ . This impression is justified when we use Stirling’s formula to approximate the binomial coefficient  $\binom{2\beta}{\beta}$  within our result. We find that the average Horton-Strahler number of  $\mathcal{C}$ -tries is asymptotically given by

$$\left(2 + \frac{1}{\beta - 1}\right) \frac{\log(\alpha) (\alpha + 1) (\alpha + 2) \sqrt{\pi}}{8 \cdot 4^\alpha \alpha^{\frac{3}{2}} \log(2)} \binom{2\alpha + 1}{\alpha - 1}.$$

Thus, only for very sparse  $\mathcal{C}$ -tries, *i.e.*  $\mathcal{C}$ -tries with few white leaves only, we have an influence of  $\beta$  on the average Horton-Strahler number. But for every fixed  $\rho$  and  $\alpha \rightarrow \infty$  also  $\beta$  tends to infinity and thus  $\frac{1}{\beta - 1}$  becomes zero. Thus it becomes possible to express the average Horton-Strahler number of  $(\alpha, \beta)$ -tries on

dependence of  $\alpha$  only. We conclude this discussion by noting that

$$\lim_{\alpha \rightarrow \infty} \frac{(\alpha + 1)(\alpha + 2)\sqrt{\pi}}{4^{\alpha+1}\alpha^{\frac{3}{2}}} \binom{2\alpha + 1}{\alpha - 1} = \frac{1}{2}$$

holds. Thus we have the following corollary:

**Corollary 1.** *Under the assumption that the number of internal nodes  $\alpha$  and the number of white leaves  $\beta$  grow in some fixed proportion the average Horton-Strahler number of  $(\alpha, \beta)$ -tries is asymptotically given by*

$$\frac{\log(\alpha)}{2 \log(2)}.$$

**Remark.** We can also conclude the result of the previous corollary by using the multivariate Darboux-method presented in [5] in order to approximate the coefficient of the leading term in (7) and the number of  $\mathcal{C}$ -tries of size  $(\alpha, \beta)$ . In that case, because of side-conditions given by the method,  $\rho = \frac{\alpha}{\beta}$  has to be strictly larger than one. However, it is impossible to derive the more accurate results presented in Lemma 3 and Theorem 1 in that way.

We will now use our generating functions to derive an asymptotic equivalent for the average Horton-Strahler number for  $\mathcal{C}$ -tries of size  $\alpha$ . As methodology is much more developed for univariate generating functions it is possible to derive results of higher precision. Note, that in the uniform model it is not possible to derive a univariate result with respect to the number of white leaves since all the generating functions would count infinitely many  $\mathcal{C}$ -tries of any given size  $\beta$ . This is due to the fact, that the number of white leaves does not limit the number of internal nodes even if we fix the Horton-Strahler number of the  $\mathcal{C}$ -tries considered. Thus, we return to (6) and set  $y = 1$  in all parts of the generating function. It is obvious that we find the same integral as a representation of  $\sum_{n \geq 1} (v_2(n) + 1)u^n$  as in the bivariate case since the substitution  $u = \exp(-t)$  yields the same result even if we set  $y$  to 1. But now it makes sense to consider terms of lower significance also, because the  $\mathcal{O}$ -transfer method for univariate generating functions makes it possible to translate them into the right contributions for the asymptotics in question. Therefore we sum the residues of the singularities at  $s \in \{1, -n, \chi_k\}$ ,  $n \in \mathbb{N} \cup \{0\}$ ,  $k \in \mathbb{Z} \setminus \{0\}$ , and multiply them by the expansion of the factor  $\frac{y\sqrt{u(1-u)}}{u} \Big|_{y=1}$  which gives us

$$2 + \left( \frac{2 \ln(t) + 2\gamma - 2 \ln(\pi) - 3 \ln(2)}{4 \ln(2)} \right) t + \sum_{k \neq 0} \frac{\Gamma(\chi_k)\zeta(\chi_k)}{\ln(2)} t^{1-\chi_k} + \mathcal{O}(t^2). \quad (9)$$

Again, we are interested in an expansion at the dominant singularity which is  $x = \frac{1}{2+4y} \Big|_{y=1} = \frac{1}{6}$ . Thus we have to set  $t = 2\sqrt{6}\sqrt{1-6x}$  in order to resubstitute

$t$  in (9) which yields the desired expansion:

$$\begin{aligned} & -\frac{\sqrt{6}\sqrt{1-6x}\ln((1-6x)^{-1})}{2\ln(2)} + \frac{\sqrt{6}\sqrt{1-6x}(2\gamma - 2\ln(\pi) + \ln(3))}{2\ln(2)} \\ & + \sum_{k \neq 0} \frac{\Gamma(\chi_k)\zeta(\chi_k)}{\ln(2)} (2\sqrt{6})^{1-\chi_k} (1-6x)^{\frac{1-\chi_k}{2}} + \mathcal{O}(|1-6x|) \\ = & -\frac{\sqrt{6}\sqrt{1-6x}\ln((1-6x)^{-1})}{2\ln(2)} + \frac{\sqrt{6}\sqrt{1-6x}(2\gamma - 2\ln(\pi) + \ln(3))}{2\ln(2)} \\ & + \frac{2\sqrt{6}}{\ln(2)} \sum_{k \neq 0} \Gamma(\chi_k)\zeta(\chi_k)e^{-\log_2(6)\pi ik} (1-6x)^{\frac{1-\chi_k}{2}} + \mathcal{O}(|1-6x|). \end{aligned}$$

Now we can apply the transfer formulae according to [8]. We have

$$[z^n](1-z)^\alpha \sim \frac{n^{-\alpha-1}}{\Gamma(-\alpha)} + \mathcal{O}(n^{-\alpha-2})$$

and

$$[z^n](1-z)^{\frac{1}{2}} \ln((1-z)^{-1}) \sim -\frac{1}{\sqrt{\pi n^3}} \left( \frac{1}{2} \ln(n) + \frac{\gamma + 2 \log(2) - 2}{2} + \mathcal{O}\left(\frac{\ln(n)}{n}\right) \right)$$

which provide the following lemma:

**Lemma 4.** *The coefficient of  $M(x, 1)$  at  $x^\alpha$  is asymptotically given by*

$$\begin{aligned} & \frac{6^{\frac{1}{2}+\alpha}(\ln(\alpha) - \gamma + 2\ln(2) - 2 + 2\ln(\pi) - \ln(3))}{4 \cdot \alpha^{\frac{3}{2}} \sqrt{\pi} \ln(2)} \\ & + \frac{2\sqrt{6}}{\ln(2)} \sum_{k \neq 0} \Gamma(\chi_k)\zeta(\chi_k)e^{-\log_2(6)\pi ik} 6^\alpha \alpha^{\frac{\chi_k}{2} - \frac{3}{2}} / \Gamma\left(\frac{\chi_k}{2} - \frac{1}{2}\right) + \mathcal{O}\left(\alpha^{-\frac{5}{2}}\right). \end{aligned}$$

Finally, this quantity has to be divided by the asymptotic number of  $\alpha$ -tries which is known (see [20]) to be given by  $6^{\frac{1}{2}+\alpha}\alpha^{-\frac{3}{2}}/(2\sqrt{\pi}) + \mathcal{O}(\alpha^{-\frac{5}{2}})$ . After numerous simplifications we find:

**Theorem 2.** *On the assumption that all  $\mathcal{C}$ -tries of the same size are equally likely the average Horton-Strahler number of an  $\alpha$ -trie is asymptotically given by*

$$\frac{1}{2} \log_2 \left( \frac{4}{3} \pi^2 \alpha \right) - \frac{\gamma + 2}{2\ln(2)} + \Delta \left( \log_2 \left( \frac{\alpha}{6} \right) \right) + \mathcal{O}(\alpha^{-1}),$$

$\chi_k = \frac{2\pi ik}{\ln(2)}$ ,  $\alpha \rightarrow \infty$ . The function  $\Delta(x)$  is a periodic function of small modulus ( $|\Delta(x)| < 0.041$ ) and possesses the following representation as a Fourier series:

$$\Delta(x) := \frac{1}{\ln(2)} \sum_{k \neq 0} (\chi_k - 1) \Gamma\left(\frac{\chi_k}{2}\right) \zeta(\chi_k) e^{\pi i k x}.$$

**Remark.** Note that this is exactly the same result as for non-colored extended binary trees given in (1) when setting  $\alpha$  to  $6\alpha$ . The same effect with a different constant can be observed for the average stack-number. In that case we have to set  $\alpha$  to  $\frac{2}{3}\alpha$  in order to get the same leading term as for non-colored trees. The fact, that there are different constants for different parameters, supports a conjecture stated in [20] which says that it seems to be impossible to conclude the behavior of  $\mathcal{C}$ -tries with respect to “traversal-parameters” from the well know results for ordinary extended binary trees (*e.g.* by a simple rearrangement together with an appropriate weighting of the trees). The bound  $|\Delta(x)| < 0.041$  can be found by means of numerical studies.

### 3. VISUALIZATION AND CONCLUSIONS

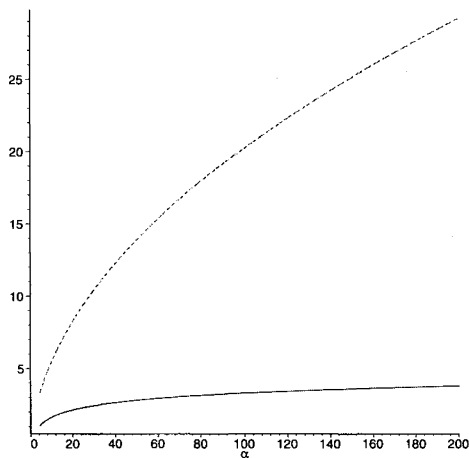


FIGURE 4. The average stack-number (upper graph) and Horton-Strahler number (lower graph) on dependence of the number of internal nodes  $\alpha$ .

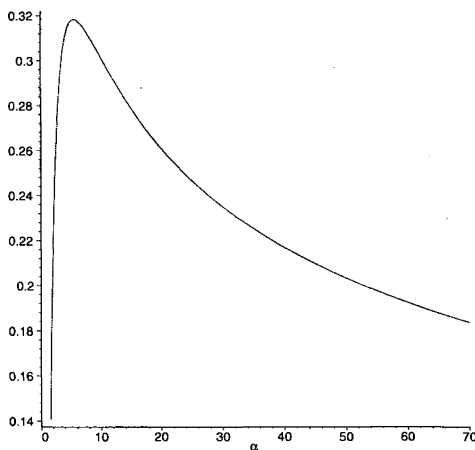


FIGURE 5. The ratio of the average Horton-Strahler number and the average stack-number.

In this section we will provide some plots of the results presented in [20], in [21] and in this paper. This is how we are going to compare the average stack-size with



the average Horton-Strahler number which are related in the following way: if we think of applications of the Horton-Strahler number such as a tree traversal or the evaluation of an arithmetic expression, the stack-size of the corresponding tree describes the amount of space needed when we apply a usual preorder traversal or a simple traversal strategy for evaluation. Those methods can be optimized with respect to the amount of space needed by easing the restriction that subtrees must be visited in a fixed order. The space requirement of the resulting strategy is described by the Horton-Strahler number. Therefore, we will speak of an *economy of space* when comparing both parameters.

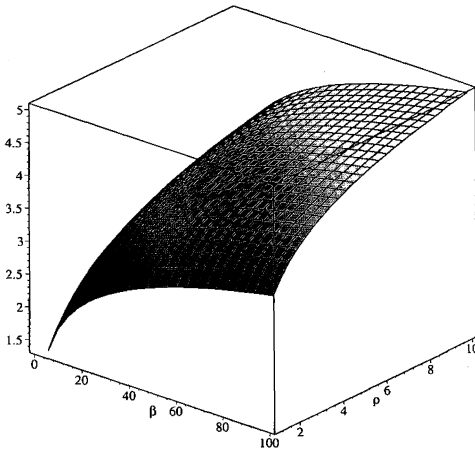


FIGURE 6. A plot of the average Horton-Strahler number on dependence of  $\rho$  and  $\beta$ .

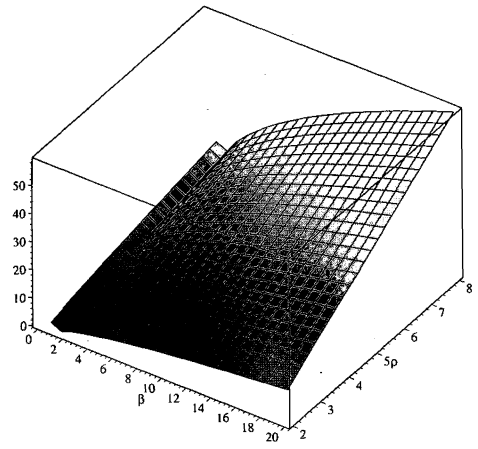


FIGURE 7. The difference of the average stack-number and the average Horton-Strahler number on dependence of  $\rho$  and  $\beta$ .

The first plot is presented in Figure 4. It shows the absolute values of the average stack-number and the average Horton-Strahler number. As we can see the order of growth of both graphs is quite different. Even if the total stack-number for  $\mathcal{C}$ -tries of size  $\alpha$  is small, the relative economy of space implied by the application of the optimized algorithms related to the Horton-Strahler number is remarkable (as we can see in Fig. 5). In the bivariate setting a similar behavior can be found. If we take a look at Figure 6 we see that the average Horton-Strahler number for  $\mathcal{C}$ -tries of size  $(\alpha, \beta)$  grows slowly and is of small value even for *sparse*  $\mathcal{C}$ -tries, *i.e.* for  $\mathcal{C}$ -tries with a large internal structure but only a few white leaves. As we would have expected from the univariate case, Figure 7 shows that the economy of space grows the larger the  $(\alpha, \beta)$ -tries become. We can also observe that the advantage of the optimized algorithms gets larger with  $\rho$  growing and not only the relative but also the total economy of space get large when the  $\mathcal{C}$ -tries become sparse. For

example on the assumption of  $\rho = 8$  the total economy of space is about 60 for  $(\alpha, \beta)$ -tries with only 20 white leaves.

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